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Discrete-Time Negative Imaginary Systems[★]

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Abstract

In this paper we introduce the notion of a discrete-time negative imaginary system and we investigate its relations with discrete-time positive real system theory. In the framework presented here, discrete-time negative imaginary systems are defined in terms of a sign condition that must be satisfied in a domain of analyticity of the transfer function, in analogy with the case of discrete-time positive real functions, as well as analogously to the continuous-time case. This means in particular that we do not need to restrict our notions and definitions to systems with rational transfer functions. We also provide a discrete-time counterpart of the different notions that have appeared so far in the literature within the framework of strictly positive real and in the more recent theory of strictly negative imaginary systems, and to show how these notions are characterized and linked to each other. Stability analysis results for the feedback interconnection of discrete-time negative imaginary systems are also derived.

Key words: Discrete-time negative imaginary systems, discrete-time negative imaginary lemma, discrete-time positive real systems.

1 Introduction

The theory of positive real (PR) systems dates back to the early 1930s [5], and is regarded as one of the cornerstones of systems and control theory, and in particular of passivity theory. For a summary of the historic and recent contributions in this area, we refer the reader to the important monographs [1], [4]. A promising new development in the area has been the introduction of the notion of *negative imaginary (NI) systems*, see [12,29,15] and the references therein. The definition of negative imaginary systems imposes a weaker restriction on the relative degree of the transfer function with respect to the one for positive real systems, and does not prohibit the case of all unstable transmission zeros. In the past few years, a rich stream of literature flourished on negative imaginary systems, including extensions

to infinite dimensional systems [18], Hamiltonian systems [25], descriptor systems [14], lossless negative imaginary systems [30] and mixtures of negative imaginary and small-gain properties [19] to mention only a few. The theory developed in these contributions has been proved to be useful in a range of applications including modelling and control of undamped or lightly damped flexible structures with colocated position sensors and force actuators [20,3], nanopositioning control due to piezoelectric transducers and capacitive sensors (e.g. [2,17,16]) and multi-agent networked systems (e.g. [6,26,27]). This theory provides a very general technique for finding an appropriate state-feedback controller and this is particularly useful when the underlying model is repeatedly derived by system identification techniques (see e.g. [31,33,32] and references therein).

An important gap in the current literature – that the present paper attempts to fill – is the lack of a definition of negative imaginary (and strictly negative imaginary) function for discrete-time systems. Furthermore, so far [9] and [8] have been the only contributions which attempted to address the general case of a definition of negative imaginary system for non necessarily rational transfer functions and then recover, in the symmetric rational case, the standard definition given in the foundational paper [12].

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The main contribution of this paper is to introduce the notion of discrete-time negative imaginary systems for the first time. This definition is given in the general non-rational setting and then is specialized for rational transfer functions, and expressed in terms of a sign constraint on the unit circle. We also introduce different notions of *strictly negative imaginary* discrete-time transfer functions that parallel the continuous-time definitions given so far. Finally, the relations between discrete-time and continuous-time negative imaginary systems are elucidated. We also provide a discrete-time negative imaginary lemma which yields a complete state-space characterization of discrete-time negative imaginary systems and a stability analysis result for the feedback interconnection of discrete-time negative imaginary systems.

Notice that negative imaginary system theory has already been proven to be very useful in the continuous-time; hence developing a discrete-time counterpart of this theory is particularly significant and promising in view of the pervasive role of digital control in modern applications.

Notation. Given a matrix A , the symbol A^\top denotes the transpose of A and A^* denotes the complex conjugate transpose of A . We denote by $\sigma(A)$ the set of singular values of the matrix A and by $\underline{\sigma}(A)$ the smallest of such singular values. The usual notations of ≥ 0 and > 0 are used to denote positive semidefiniteness and positive definiteness of Hermitian matrices, respectively. Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be analytic or harmonic in a certain region Ω of \mathbb{C} , then G is said to have full normal rank if there exists $z \in \Omega$ such that $\det[G(z)] \neq 0$.

2 Discrete-Time Positive Real Functions

In this section, for the sake of completeness we briefly recall the most important notions and results of discrete-time positive real systems. The definition of *discrete-time positive real* function was introduced for the first time in the literature by Hitz and Anderson in [10], and is recalled below.

Definition 2.1 [10]. *The function $F : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ is discrete-time positive real (PR) if*

- $F(z)$ is analytic in $\{z \in \mathbb{C} : |z| > 1\}$;
- $F(z)$ is real when z is real and positive;
- $F(z)^* + F(z) \geq 0$ for all $|z| > 1$.

Similarly to what happens in the continuous-time for rational functions, discrete-time positive realness can be characterized in terms of conditions involving properties of the restriction of the matrix function to the unit circle.

Theorem 2.1 [10, LEMMA 2]. *Let $F : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper transfer function. Then, $F(z)$ is PR if and only if*

- $F(z)$ has no poles in $\{z \in \mathbb{C} : |z| > 1\}$;
- $F(e^{i\theta})^* + F(e^{i\theta}) \geq 0$ for all $\theta \in [0, 2\pi)$ except for the values of θ for which $z = e^{i\theta}$ is a pole of $F(z)$;

- If $z_0 = e^{i\theta_0}$, with $\theta_0 \in [0, 2\pi)$, is a pole of $F(z)$, then it is a simple pole and the normalized residual matrix

$$K_0 \stackrel{\text{def}}{=} \frac{1}{z_0} \lim_{z \rightarrow z_0} (z - z_0) F(z)$$

is Hermitian and positive semidefinite.

We now present a definition of discrete-time strictly positive real systems. We warn the reader that many different definitions have been proposed for this concept that can indeed be distinguished via several grades of strength, see e.g. [4,11]. In this paper, we shall only need two of such grades, that will be referred to as *strongly* and *weakly* strictly positive realness.

Definition 2.2 *Let $F : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, proper transfer function. Then, $F(z)$ is discrete-time strongly strictly positive real (SSPR) if for some $\delta \in (0, 1)$, the transfer function $F(\delta z)$ is PR and $F(z) + F(1/z)^\top$ has full normal rank.*

The following result shows that in the case of rational functions the property of SSPR is equivalent to an analyticity condition and a sign condition restricted to the unit circle.

Theorem 2.2 *Let $F : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper transfer function. Then, $F(z)$ is SSPR if and only if*

- $F(z)$ has all its poles in a disk of radius $\rho \in [0, 1)$;
- $F(e^{i\theta}) + F(e^{i\theta})^* > 0$ for all $\theta \in [0, 2\pi)$.

Proof: Necessity of the first condition is obvious. Necessity of the second immediately follows from the fact the unit circle is in the interior of the domain of analyticity and by the full normal rank assumption. As for sufficiency, since the unit circle is closed, condition $F(e^{i\theta}) + F(e^{i\theta})^* > 0$ for all $\theta \in [0, 2\pi)$ implies coercivity, i.e. there exists $\sigma_0 > 0$ such that $F(e^{i\theta}) + F(e^{i\theta})^* > \sigma_0 I$ for all $\theta \in [0, 2\pi)$. Therefore, there exists $\rho \in [0, 1)$ such that $F(\rho e^{i\theta}) + F(\rho e^{i\theta})^* > 0$ for all $\theta \in [0, 2\pi)$, so that $F_1(z) \stackrel{\text{def}}{=} F(\rho z)$ is PR. ■

Remark 2.1 The conditions of Theorem 2.2 are much simpler than those of its continuous-time counterpart (see e.g. [4, Theorem 2.47] and [11, Lemma 6.1]) because positivity on the unit circle \mathbb{T} implies coercivity in view of the closure of the unit circle \mathbb{T} (as opposed to the fact that the imaginary axis is not closed). As we shall see later, this is not the case for discrete-time negative imaginary (NI) systems for which the relevant boundary curve is the intersection between \mathbb{T} and the open upper half complex plane. Therefore, the relevant boundary curve is not closed as the zero and infinity discrete frequencies are not in this curve. This fact complicates the derivations and the results in the NI case.

The next result is the discrete-time counterpart of the so-called *positive real lemma*, a cornerstone of modern control theory that has generated an endless stream of literature.

Lemma 2.1 [DISCRETE-TIME POSITIVE REAL LEMMA, [10, LEMMA 3]]. Let $F : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper transfer function with no poles in $|z| >$

1 and simple poles only on $|z| = 1$. Let $\begin{bmatrix} A & B \\ C & D \end{bmatrix}$ be a minimal

realization of $F(z)$. Then $F(z)$ is discrete-time positive real if and only if there exist a real matrix $X = X^\top > 0$ and real matrices L and W such that

$$X - A^\top X A = L^\top L, \quad (1)$$

$$C^\top - A^\top X B = L^\top W, \quad (2)$$

$$D^\top + D - B^\top X B = W^\top W. \quad (3)$$

3 Discrete-Time Negative Imaginary Functions

We now introduce the following standing assumption that will be used throughout the rest of the paper.

Assumption 3.1 We henceforth restrict our attention to only symmetric transfer functions.

As discussed in [9], the case of symmetric transfer function is the most important one, because it encompasses both the scalar case, and the case of a transfer function of a reciprocal m -port electrical network.¹ To the best of our knowledge, all the negative imaginary transfer functions considered or studied in the literature so far are symmetric (see e.g. the transfer functions from a force actuator to a corresponding collocated position sensor — for instance, a piezoelectric sensor — in a lightly damped or undamped structure), even though the real, rational definitions of negative imaginary systems in [12,29,15] allow for non-symmetric transfer functions.

We now present a definition of negative imaginary functions in the discrete-time case.

Definition 3.1 Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real transfer function. We say that $G(z)$ is discrete-time negative imaginary (NI) if

- (i) $G(z)$ is analytic in $\{z \in \mathbb{C} : |z| > 1\}$;
- (ii) $i [G(z) - G(z)^*] \geq 0$ for all $z \in \mathbb{C}$ such that $|z| > 1$ and $\Im m(z) > 0$;
- (iii) $i [G(z) - G(z)^*] = 0$ for all $z \in \mathbb{C}$ such that $|z| > 1$ and $\Im m(z) = 0$;
- (iv) $i [G(z) - G(z)^*] \leq 0$ for all $z \in \mathbb{C}$ such that $|z| > 1$ and $\Im m(z) < 0$.

¹ We recall that the only way to obtain a non-symmetric transfer function of an m -port electrical network is to employ gyrators, whose physical implementation requires the use of active components but that cannot be physically implemented with arbitrary precision.

The conditions (ii)-(iv) in Definition 3.1 are a *skew imaginary condition* on the open set $\Omega = \{z \in \mathbb{C} : |z| > 1\}$.

Remark 3.1 Note that if the real transfer function $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ satisfies the conditions in Definition 3.1, then $G(z)$ is symmetric, i.e., $G(z) = G(z)^\top$ for all $z \in \mathbb{C}$ such that $|z| > 1$. This can be seen as follows: since $G(z)$ is real, if $z \in \mathbb{R}$ then $G(z) \in \mathbb{R}$. Let $z \in \mathbb{R}$ and $|z| > 1$. From (iii), we get $G(z) = G(z)^\top$. Therefore, each entry $\Delta_{ij}(z)$ of the matrix valued function $\Delta(z) \stackrel{\text{def}}{=} G(z) - G(z)^\top$ is a function which is analytic in $\{z \in \mathbb{C} : |z| > 1\}$ and is zero for any real z in the domain of analyticity. Then, in view of the principle of identity of analytic functions (see e.g. Corollary to Theorem 10.18 in [22, page 209]) $\Delta_{ij}(z) = 0$ in the entire domain of analyticity so that $G(z) = G(z)^\top$ for all z in the domain of analyticity, i.e., $|z| > 1$.

Conditions (iii)-(iv) in Definition 3.1 are redundant in the rational case, as the following result establishes.

Lemma 3.1 Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational transfer function. If $G(z)$ satisfies (i)-(ii) of Definition 3.1, then it also satisfies (iii)-(iv).

Proof: If $G(z)$ satisfies (ii), then $i [G(z)^\top - G(\bar{z})] \geq 0$ for all $z \in \mathbb{C}$ such that $|z| > 1$ and $\Im m(z) > 0$, since $G(z)^* = G(\bar{z})^\top$. Defining $w \stackrel{\text{def}}{=} \bar{z}$, such condition can be re-written as $i [G(w)^* - G(w)] \geq 0$ for all $w \in \mathbb{C}$ such that $|w| > 1$ and $\Im m(w) < 0$, which is exactly (iv) of Definition 3.1. Finally, since (ii) and (iv) hold, then (iii) must also hold by continuity. ■

We now prove the counterpart of Theorem 2.1 for the case of discrete-time symmetric negative imaginary functions. This result provides a characterization of rational NI systems in terms of a domain of analyticity and conditions referred to the unit circle. First, however, we recall that given a real rational function $G(z)$ and a simple pole $p \in \mathbb{C}$ of $G(z)$, we have a unique decomposition $G(z) = G_1(z) + A/(z-p)$, where $G_1(z)$ is a rational function which is analytic in an open set containing p and the (non-zero) matrix A is the residue corresponding to the pole p . If p is a double pole of $G(z)$, we have the unique decomposition $G(z) = G_1(z) + A_1/(z-p) + A_2/(z-p)^2$, where the matrix A_1 is the residue corresponding to the pole p . In this case, by analogy, we define the (non-zero) matrix A_2 to be the *quadratic residue* corresponding to the pole p . If $G(z)$ has a pole at infinity, it can be uniquely decomposed as $G(z) = G_1(z) + P(z)$, where $G_1(z)$ is a rational proper function and $P(z) = \sum_{i=1}^k A_i z^i$ is a homogeneous polynomial in z . We refer to A_i as the i -th coefficient in the expansion at infinity of $G(z)$.

Lemma 3.2 Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper transfer function. Then, $G(z)$ is NI if and only if

- (i) $G(z)$ has no poles in $|z| > 1$;

- (ii) $i[G(e^{i\theta}) - G(e^{i\theta})^*] \geq 0$ for all $\theta \in (0, \pi)$ except for the values of θ for which $z = e^{i\theta}$ is a pole of $G(z)$;
- (iii) if $z_0 = e^{i\theta_0}$, with $\theta_0 \in (0, \pi)$, is a pole of $G(z)$, then it is a simple pole and the normalized residual matrix

$$K_0 \stackrel{\text{def}}{=} \frac{1}{z_0} \lim_{z \rightarrow z_0} (z - z_0) iG(z) \quad (4)$$

is Hermitian and positive semidefinite;

- (iv) if $z_0 = 1$ is a pole of $G(z)$, then it is at most a double pole. Moreover, its residue A_1 and its quadratic residue A_2 (when the pole is simple it is assumed that $A_2 = 0$) are Hermitian matrices satisfying $A_2 \geq 0$ and $A_1 \geq A_2$;
- (v) if $z_0 = -1$ is a pole of $G(z)$, then it is at most a double pole. Moreover, its residue A_1 and its quadratic residue A_2 (when the pole is simple it is assumed that $A_2 = 0$) are Hermitian matrices satisfying $A_2 \leq 0$ and $A_1 \geq -A_2$.

Proof: The idea of the proof is the following: we introduce a bilinear transform and show that it maps continuous-time NI systems into discrete-time NI systems. Then we show that the conditions of Lemma 3.2 are mapped by the bilinear transform into the necessary and sufficient conditions derived in [9, Lemma 3.1] for a continuous-time system to be NI. Let $G(z)$ be discrete-time real, symmetric and rational, and define

$$G_c(s) \stackrel{\text{def}}{=} G\left(\frac{1+s}{1-s}\right).$$

Consider the bilinear transform

$$z = \frac{1+s}{1-s},$$

and let $z = \sigma + i\omega$. It is found that

$$s = \frac{z-1}{z+1} = \frac{\sigma^2 + \omega^2 - 1}{(\sigma+1)^2 + \omega^2} + 2i \frac{\omega}{(\sigma+1)^2 + \omega^2}. \quad (5)$$

Firstly, $G(z)$ is NI if and only if $G_c(s)$ is NI as a continuous-time transfer function. Indeed, in view of (5), $G(z)$ is analytic in $|z| > 1$ if and only if $G_c(s)$ is analytic in $\Re\{s\} > 0$. The rest of the proof of this part follows directly from the definitions, using the fact that $\Im\{z\} > 0$ (resp. $\Im\{z\} < 0$ and $\Im\{z\} = 0$) is equivalent to $\omega > 0$ (resp. $\omega < 0$ and $\omega = 0$), which in turn is equivalent to $\Im\{s\} > 0$ (resp. $\Im\{s\} < 0$ and $\Im\{s\} = 0$).

Secondly, the following facts are easy to check:

- $G(z)$ has no poles in $|z| > 1$ if and only if $G_c(s)$ has no poles in $\Re\{s\} > 0$;
- Let $z_0 \stackrel{\text{def}}{=} e^{i\theta_0}$ with $\theta_0 \in (0, \pi)$. Using (5) we see that

$$s_0 \stackrel{\text{def}}{=} \frac{z_0 - 1}{z_0 + 1} = \frac{e^{i\theta_0} - 1}{e^{i\theta_0} + 1} = i \frac{\sin \theta_0}{1 + \cos \theta_0},$$

which shows that z_0 is a pole of $G(z)$ if and only if $i\omega_0$, with $\omega_0 \stackrel{\text{def}}{=} \frac{\sin \theta_0}{1 + \cos \theta_0} > 0$, is a purely imaginary pole of $G_c(s)$. Moreover, $i[G(e^{i\theta}) - G(e^{i\theta})^*] \geq 0$ for all $\theta \in (0, \pi)$ such that $e^{i\theta}$ is not a pole of $G(z)$ if and only if $i[G_c(i\omega) - G_c(i\omega)^*] \geq 0$ for all $\omega \in (0, \infty)$ such that $i\omega$ is not a pole of $G_c(s)$;

- Let $z_0 \stackrel{\text{def}}{=} e^{i\theta_0}$ with $\theta_0 \in (0, \pi)$. Then z_0 , with $\theta_0 \in (0, \pi)$, is a pole of G if and only if $i\omega_0$, with $\omega_0 \stackrel{\text{def}}{=} \frac{\sin \theta_0}{1 + \cos \theta_0} > 0$, is a purely imaginary pole of G_c . Moreover, they are poles with the same multiplicity. Finally, z_0 is a simple pole of $G(z)$ with residue being the matrix K if and only if $i\omega_0$ is a simple pole of $G_c(s)$ with residue being the matrix $H \stackrel{\text{def}}{=} \frac{e^{-i\theta_0}}{1 + \cos \theta_0} K$. Notice that the normalized residual matrix K_0 of $G(z)$, as defined in (4), is positive semi-definite if and only if $\frac{i}{z_0} K$ is positive semi-definite and, hence, if and only if iH is positive semi-definite.
- $z_0 = 1$ is a pole of $G(z)$ if and only if $s_0 = 0$ is a pole of G_c . This fact follows straightforwardly from (5). Moreover, they are poles with the same multiplicity. If this multiplicity is strictly greater than 2, then $G(z)$ is trivially not NI. If this multiplicity is at most 2, then the residue A_{s1} and the quadratic residue A_{s2} corresponding to s_0 are related to the residue A_1 and the quadratic residue A_2 corresponding to z_0 by: $A_{s2} = \frac{1}{4}A_2$ and $A_{s1} = \frac{1}{2}(A_1 - A_2)$, since

$$G(z) = G_1(z) + \frac{A_1}{z-1} + \frac{A_2}{(z-1)^2},$$

where $G_1(z)$ is analytic in an open set containing $z_0 = 1$, and

$$\begin{aligned} G_c(s) &= G_{c,1}(s) + \frac{A_1}{\frac{1+s}{1-s} - 1} + \frac{A_2}{\left(\frac{1+s}{1-s} - 1\right)^2} \\ &= \left(G_{c,1}(s) - \frac{A_1}{2} + \frac{A_2}{4}\right) + \frac{A_1 - A_2}{2s} + \frac{A_2}{4s^2} \\ &= \left(G_{c,1}(s) - \frac{A_1}{2} + \frac{A_2}{4}\right) + \frac{A_{s1}}{s} + \frac{A_{s2}}{s^2}, \end{aligned}$$

where $G_{c,1}(s) - \frac{A_1}{2} + \frac{A_2}{4}$ is analytic in an open set containing $s_0 = 0$.

- $z_0 = -1$ is a pole of $G(z)$ if and only if ∞ is a pole of G_c . Moreover, they are poles with the same multiplicity. If this multiplicity is strictly greater than 2, then $G(z)$ is trivially not NI. In the case in which this multiplicity is at most 2, the first coefficient A_{s1} and the second coefficient A_{s2} in the expansion at infinity of $G_c(s)$ are connected to the residue A_1 and the quadratic residue A_2 corresponding to z_0 by: $A_{s2} = \frac{1}{4}A_2$ and $A_{s1} = -\frac{1}{2}(A_1 + A_2)$ since

$$G(z) = G_1(z) + \frac{A_1}{z+1} + \frac{A_2}{(z+1)^2},$$

where $G_1(z)$ is analytic in an open set containing $z_0 = -1$,

and

$$\begin{aligned} G_c(s) &= G_{c,1}(s) + \frac{A_1}{\frac{1+s}{1-s} + 1} + \frac{A_2}{\left(\frac{1+s}{1-s} + 1\right)^2} \\ &= \left(G_{c,1}(s) + \frac{A_1}{2} + \frac{A_2}{4}\right) - \frac{A_1 + A_2}{2}s + \frac{A_2}{4}s^2 \\ &= \left(G_{c,1}(s) + \frac{A_1}{2} + \frac{A_2}{4}\right) + A_{1s}s + A_{2s}s^2, \end{aligned}$$

where $G_{c,1}(s) + \frac{A_1}{2} + \frac{A_2}{4}$ is rational and proper.

Now, we apply [9, Lemma 3.1] in both directions and get the desired result. ■

Remark 3.2 In Definition 3.1 we need to assume symmetry of the transfer function matrix in order to introduce the notion of a NI system as a property that is defined in the domain of analyticity: this definition is the analogue to the classic definition of PR systems and has the important advantage of considering a general setting that does not require rationality assumptions. Note, however, that if one is only interested in the rational case, it is possible to consider conditions (i)-(v) of Lemma 3.2 as the definition of rational NI transfer functions and this clearly does not require any symmetry assumption. This is indeed the route taken in the first papers on continuous-time NI systems, see [12,20]. A similar observation can be made for the definition of strictly negative imaginary systems given below.

The reader can check that, as long as one considers only rational transfer functions, all the results derived in this paper can be generalized to the case of non-symmetric transfer functions.

We now define the notions of strictly negative imaginary systems in discrete-time.

Definition 3.2 Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real transfer function. Then, $G(z)$ is discrete-time strongly strictly negative imaginary (SSNI) if for some $\delta \in (0, 1)$, the transfer function $G(\delta z)$ is NI and $i[G(z) - G(1/z)^\top]$ has full normal rank.

Now, we show that SSNI as defined in Definition 3.2 can be equivalently checked via conditions on the domain of analyticity.

Lemma 3.3 Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real transfer function. Then, $G(z)$ is SSNI if and only if there exists $\delta \in (0, 1)$ such that

- (i) $G(z)$ is analytic in $\{z \in \mathbb{C} : |z| > \delta\}$;
- (ii) $i[G(z) - G(z)^*] > 0$ for all $z \in \mathbb{C}$ such that $|z| > \delta$ and $\Im\{z\} > 0$;
- (iii) $i[G(z) - G(z)^*] = 0$ for all $z \in \mathbb{C}$ such that $|z| > \delta$ and $\Im\{z\} = 0$;

- (iv) $i[G(z) - G(z)^*] < 0$ for all $z \in \mathbb{C}$ such that $|z| > \delta$ and $\Im\{z\} < 0$.

Proof: Definition 3.2 trivially gives equivalence to the existence of $\delta \in (0, 1)$ such that conditions (i)-(iv) are satisfied with non-strict inequalities in (ii) and (iv) on $i[G(z) - G(z)^*]$. We hence only need to show that the fact that G is SSNI implies that the inequalities in (ii) and (iv) are indeed strict. We prove only that (ii) is strict since (iv) follows by symmetry. Let G be analytic in $\mathbb{C}_\delta \stackrel{\text{def}}{=} \{z \in \mathbb{C} : |z| > \delta\}$ and assume by contradiction that there exist $z_0 \in \{z \in \mathbb{C} : |z| > \delta \text{ and } \Im\{z\} > 0\}$ and a nonzero vector v such that $v^*(i[G(z_0) - G(z_0)^*])v = 0$. Since G is analytic in \mathbb{C}_δ , the function $h(z) \stackrel{\text{def}}{=} v^*(i[G(z) - G(z)^*])v$ is harmonic in the same domain. Consider a real number $M > 1$ such that $M > |z_0|$ and a real number δ_1 such that $\delta < \delta_1 < |z_0|$, so that z_0 is in the interior of the compact set $\mathcal{C} \stackrel{\text{def}}{=} \{z \in \mathbb{C} : \delta_1 \leq |z| \leq M, \Im\{z\} \geq 0\}$, which is contained in \mathbb{C}_δ . Since $h(z)$ is non-negative in \mathcal{C} and $h(z_0) = 0$, then $h(z)$ restricted to \mathcal{C} attains its minimum at a point z_0 in the interior of \mathcal{C} . Hence, $h(z)$ is identically zero in \mathcal{C} and hence, in particular, it is identically zero in the (upper half of the unit circle and, by symmetry, in the whole of the) unit circle. This is a contradiction, since $i[G(z) - G(1/z)^\top]$ — that coincides with $h(z)$ in the unit circle — is required to have full normal rank by Definition 3.2. ■

We now specialize Lemma 3.3 to the unit disc. However, first we need a preliminary lemma.

Lemma 3.4 Let $g : \mathbb{C} \rightarrow \mathbb{C}$ be a scalar discrete-time, real, rational, proper transfer function. Assume that $g(z)$ is a SSNI function. If $g(1) = 0$ then the multiplicity of the zero in 1 of $g(z)$ is equal to 1. Similarly, if $g(-1) = 0$ then the multiplicity of the zero in -1 of $g(z)$ is equal to 1.

Proof: Since $g(z)$ is a SSNI function, it has no poles in 1 and we can expand $g(z)$ at 1 as

$$g(z) = \sum_{k=h}^{\infty} r_k (z-1)^k,$$

where h is the multiplicity of the zero in 1 of $g(z)$, $r_h \neq 0$, and $\limsup_{k \rightarrow \infty} |r_k|^{1/k}$ is finite so that there exists a constant c such that $|r_k|^{1/k} < c$ for all $k \geq h$. Let $z = 1 + \varepsilon e^{i\theta}$, with $\varepsilon > 0$ and $0 < \theta < \pi$. We have

$$p(z) \stackrel{\text{def}}{=} i[g(z) - g(z)^*] = \varepsilon^h [-2r_h \sin(h\theta) + \delta]$$

with $\delta \stackrel{\text{def}}{=} \varepsilon^{-h} \sum_{k=h+1}^{\infty} -2r_k \varepsilon^k \sin(k\theta)$. For ε such that $|c\varepsilon| < 1$, we have

$$|\delta| \leq \sum_{k=h+1}^{\infty} |2r_k \varepsilon^{k-h} \sin(k\theta)| \leq \varepsilon \frac{2|c|^{h+1}}{1 - c\varepsilon}$$

which is arbitrarily small for a sufficiently small ε . We can then choose ε such that $|\delta| < |2r_h|$. Assume, by contradiction, that $h > 1$ and define $\theta_1 \stackrel{\text{def}}{=} \pi/(2h)$ and $\theta_2 \stackrel{\text{def}}{=} 3\pi/(2h)$ (notice that, if $h > 1$, both θ_1 and θ_2 are in $(0, \pi)$). Then we have that $p(\varepsilon e^{i\theta_1}) = \varepsilon^h[-2r_h + \delta]$ and $p(\varepsilon e^{i\theta_2}) = \varepsilon^h[2r_h + \delta]$ have opposite signs which is a contradiction because $g(z)$ is a SSNI function so that $p(\varepsilon e^{i\theta})$ is positive for any pair $\varepsilon > 0$ and $0 < \theta < \pi$. The proof for -1 is similar. ■

Theorem 3.1 *Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper transfer function. Then $G(z)$ is SSNI if and only if*

- (i) $G(z)$ has all its poles strictly inside the unit circle;
- (ii) $i[G(e^{i\theta}) - G(e^{i\theta})^*] > 0$ for all $\theta \in (0, \pi)$;
- (iii)

$$Q_0 \stackrel{\text{def}}{=} \lim_{\theta \rightarrow 0^+} \frac{1}{\sin \theta} i[G(e^{i\theta}) - G(e^{i\theta})^*] > 0$$

(iv)

$$Q_\pi \stackrel{\text{def}}{=} \lim_{\theta \rightarrow \pi^-} \frac{1}{\sin \theta} i[G(e^{i\theta}) - G(e^{i\theta})^*] > 0$$

Proof: Necessity of (i) and (ii) is trivial from Lemma 3.3. We now prove necessity of (iii) (necessity of (iv) is similar). Assume that G is SSNI. Then clearly the limit Q_0 defined in (iii) exists and is positive semi-definite. Assume by contradiction that Q_0 is singular and let $v \in \ker Q_0$. Let $g'(z) \stackrel{\text{def}}{=} v^\top G(z)v$ and $g(z) \stackrel{\text{def}}{=} g'(z) - g'(1)$. Clearly, $g(z)$ is a rational proper SSNI function with a zero in 1 and such that

$$\lim_{\theta \rightarrow 0^+} \frac{1}{\sin \theta} i[g(e^{i\theta}) - g(e^{i\theta})^*] = 0. \quad (6)$$

By expanding $g(z)$ around 1 as

$$g(z) = \sum_{k=h}^{\infty} r_k (z-1)^k$$

we see that (6) implies that $h > 1$, which is a contradiction in view of Lemma 3.4.

As for sufficiency, assume that $G(s)$ is real symmetric and rational and that it satisfies (i), (ii), (iii) and (iv). We now show that we can choose $\rho < 1$ in such a way that

$$i[G(\rho e^{i\theta}) - G(\rho e^{i\theta})^*] > 0 \quad \forall \theta \in (0, \pi). \quad (7)$$

In view of condition (ii), we have that for all $\pi > \theta_2 > \theta_1 > 0$ there exists $\rho < 1$ such that

$$i[G(\rho e^{i\theta}) - G(\rho e^{i\theta})^*] > 0 \quad \forall \theta \in [\theta_1, \theta_2], \quad (8)$$

so that it is sufficient to show that given an arbitrarily small θ_1 and an arbitrarily large θ_2 ,

$$i[G(\rho e^{i\theta}) - G(\rho e^{i\theta})^*] > 0 \quad \forall \theta \in (0, \theta_1) \quad (9)$$

and

$$i[G(\rho e^{i\theta}) - G(\rho e^{i\theta})^*] > 0, \quad \forall \theta \in (\theta_2, \pi). \quad (10)$$

As for (9), let $\delta \stackrel{\text{def}}{=} \rho e^{i\theta} - 1$ with $\theta \in (0, \theta_1)$ and consider the following expansion of $G(\delta)$:

$$G(\delta) = D_0 + \delta D_1 + \delta^2 D_2 + \dots$$

which clearly converges for a sufficiently small δ (if we considered a minimal realization $G(z) = C(zI - A)^{-1}B + D$, we would have $D_0 \stackrel{\text{def}}{=} D - C(I - A)^{-1}B$ and $D_i \stackrel{\text{def}}{=} -C(I - A)^{-i-1}B$ for $i > 1$). Since $G(z)$ is real symmetric by our standing assumption, we have $D_i = D_i^\top$. Moreover, $Q_0 \stackrel{\text{def}}{=} \lim_{\theta \rightarrow 0^+} (1/\sin \theta) i[G(e^{i\theta}) - G(e^{i\theta})^*] = -2D_1$, so that by (iv), we have $D_1 < 0$. A direct calculation gives

$$\begin{aligned} i[G(\rho e^{i\theta}) - G(\rho e^{i\theta})^*] &= -\rho \sin(\theta) 2D_1 \\ &\quad + i \sum_{j=2}^{\infty} [\delta^j - (\delta^*)^j] D_j. \end{aligned}$$

Now we observe that

$$i \sum_{j=2}^{\infty} [\delta^j - (\delta^*)^j] D_j = -2\rho \sin \theta \sum_{j=2}^{\infty} \sum_{k=0}^{j-1} [\delta^k (\delta^*)^{j-1-k}] D_j,$$

so that

$$\begin{aligned} \|i \sum_{j=3}^{\infty} [\delta^j - (\delta^*)^j] D_j\| &\leq 2\rho \sin \theta \sum_{j=2}^{\infty} j |\delta|^{j-1} \|D_j\| \\ &= 2\rho \sin \theta |\delta| \sum_{j=2}^{\infty} j |\delta|^{j-2} \|D_j\| \\ &\leq 2\rho \sin \theta |\delta| \sigma \end{aligned}$$

for a certain σ which remains bounded as $|\delta|$ tends to zero. Since, by choosing a sufficiently small δ we can make $-D_1 > \sigma |\delta| I$, we have (9). The proof of (10) is similar. ■

In analogy with the continuous-time case [8], we introduce the following definition of a weaker notion of strictly negative imaginary systems.

Definition 3.3 *The discrete-time, real, rational, proper transfer function $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ is discrete-time weakly strictly negative imaginary (WSNI) if it satisfies conditions (i) and (ii) of Theorem 3.1.*

The next lemma shows that the definition of WSNI characterizes properties on the outside of the unit disk too.

Lemma 3.5 *Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper transfer function. Then, $G(z)$ is WSNI if and only if there exists $\delta \in (0, 1)$ such that*

- (i) $G(z)$ is analytic in $\{z \in \mathbb{C} : |z| > \delta\}$;

- (ii) $i[G(z) - G(z)^*] > 0$ for all $z \in \mathbb{C}$ such that $|z| \geq 1$ and $\Im m(z) > 0$;
- (iii) $i[G(z) - G(z)^*] = 0$ for all $z \in \mathbb{C}$ such that $|z| \geq 1$ and $\Im m(z) = 0$;
- (iv) $i[G(z) - G(z)^*] < 0$ for all $z \in \mathbb{C}$ such that $|z| \geq 1$ and $\Im m(z) < 0$.

Proof: Sufficiency is trivial by restricting on $\{z \in \mathbb{C} : |z| = 1\}$. Necessity can be proven as follows: if G is WSNI, then (i) is satisfied and G is NI (from Lemma 3.2). If G is NI, then (ii)-(iv) in Definition 3.1 are satisfied. Strict inequalities in conditions (ii) and (iv) outside the unit circle are then obtained via an argument similar to that given in the proof of Lemma 3.3. Appending the $\{z \in \mathbb{C} : |z| = 1\}$ properties of G to the conditions (ii)-(iv) in Definition 3.1 (since G is WSNI) yields (ii)-(iv) above since G fulfils (i) above. ■

The following lemma relates the strong classes with the weak classes with the non-strict classes of negative imaginary systems.

Lemma 3.6 *The set of SSNI systems is contained in the set of WSNI systems which is in turn contained in the set of NI systems.*

Proof: Trivial from the definitions. ■

The following lemma relates a NI system with a PR system.

Lemma 3.7 *Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, symmetric, real, rational, proper transfer function with no poles at $z = -1$. Then, $G(z)$ is NI if and only if*

$$F(z) = \frac{z-1}{z+1} [G(z) - G(-1)] \quad (11)$$

is PR.

Proof: (Only if). The set of poles of $F(z)$ is contained in the set of poles of $G(z)$ (in fact, in (11) the pole in -1 of $\frac{z-1}{z+1}$ is cancelled by the zero in -1 of $[G(z) - G(-1)]$). Since $G(z)$ is a symmetric, real, rational, proper, NI transfer function, $F(z)$ is analytic in $|z| > 1$. Let $\theta_0 \in (0, \pi)$, and assume that $z = e^{i\theta_0}$ is not a pole of $G(z)$. Then, $z = e^{i\theta_0}$ is not a pole of $F(z)$, and a simple calculation gives

$$F(e^{i\theta_0}) + F(e^{i\theta_0})^* = \frac{\sin \theta_0}{1 + \cos \theta_0} i [G(e^{i\theta_0}) - G(e^{i\theta_0})^*] \geq 0$$

in view of Lemma 3.2.

Let us now assume that $z = e^{i\theta_0}$, with $\theta_0 \in (0, \pi)$, is a pole of $G(z)$. From Lemma 3.2, it is a simple pole, and from (11) it is also a simple pole of $F(z)$. We can write

$$G(z) = G_1(z) + \frac{A}{z - e^{i\theta_0}},$$

where $G_1(z)$ is a rational function which is analytic in an open set containing $z = e^{i\theta_0}$ and the matrix A is non-zero. Then,

$$\begin{aligned} K_0 &= e^{-i\theta_0} \lim_{z \rightarrow e^{i\theta_0}} (z - e^{i\theta_0}) i G(z) \\ &= e^{-i\theta_0} \lim_{z \rightarrow e^{i\theta_0}} (z - e^{i\theta_0}) i \left(G_1(z) + \frac{A}{z - e^{i\theta_0}} \right) = i e^{-i\theta_0} A \end{aligned}$$

is Hermitian and positive semidefinite. The normalized residue of $F(z)$ in $e^{i\theta_0}$ is given by

$$\begin{aligned} &e^{-i\theta_0} \lim_{z \rightarrow e^{i\theta_0}} (z - e^{i\theta_0}) F(z) \\ &= e^{-i\theta_0} \lim_{z \rightarrow e^{i\theta_0}} \frac{z-1}{z+1} \left[(z - e^{i\theta_0}) G(z) - (z - e^{i\theta_0}) G(-1) \right] \\ &= e^{-i\theta_0} \frac{e^{i\theta_0} - 1}{e^{i\theta_0} + 1} A = \frac{\sin \theta_0}{1 + \cos \theta_0} i e^{-i\theta_0} A \geq 0. \end{aligned}$$

Let us now consider the case $\theta_0 = 0$, i.e., $z_0 = e^{i\theta_0} = 1$. If $G(z)$ has no poles at $z_0 = 1$, neither does $F(z)$. In this case, $F(1) = 0$, which gives $F(1) + F(1)^* = 0 \geq 0$. If $G(z)$ has a simple pole at $z_0 = 1$, then $F(z)$ has no poles at $z_0 = 1$. In this case, $G(z) = G_1(z) + \frac{A}{z-1}$, where $G_1(z)$ is a rational function which is analytic in an open set containing $z_0 = 1$, and where $A \geq 0$ from (iv) in Lemma 3.2 (because the quadratic residual is zero). Thus,

$$F(z) = \frac{z-1}{z+1} \left[G_1(z) + \frac{A}{z-1} - G(-1) \right],$$

so that $F(1) = A/2$, and $F(1) + F(1)^* = A \geq 0$. Now, consider the case in which $G(z)$ has a double pole at $z_0 = 1$. In this case, we can write $G(z) = G_1(z) + \frac{A_1}{z-1} + \frac{A_2}{(z-1)^2}$, where $G_1(z)$ is a rational function which is analytic in an open set containing $z_0 = 1$, $A_1 \geq A_2$ and $A_2 \geq 0$. In this case,

$$\begin{aligned} F(z) &= \frac{z-1}{z+1} \left[G_1(z) + \frac{A_1}{z-1} + \frac{A_2}{(z-1)^2} - G(-1) \right] \\ &= \left[\frac{z-1}{z+1} G_1(z) + \frac{A_1}{z+1} - \frac{z-1}{z+1} G(-1) - \frac{A_2}{2(z+1)} \right] \\ &\quad + \frac{A_2}{2(z-1)}. \end{aligned}$$

Since $G_1(z)$ is analytic in an open set containing $z_0 = 1$, $\frac{z-1}{z+1} G_1(z) + \frac{A_1}{z+1} - \frac{z-1}{z+1} G(-1) - \frac{A_2}{2(z+1)}$ is also analytic in an open set containing $z_0 = 1$. Thus, $F(z)$ has a simple pole at $z_0 = 1$, and the corresponding residue $A_2/2$ is positive semidefinite (notice that in this case the residue and the normalized residue coincide because $z_0 = 1$).

Let us finally consider the case $\theta_0 = \pi$, i.e., $z_0 = e^{i\theta_0} = -1$. We know that $G(-1)$ is finite and hence $F(-1)$ is finite as well. Moreover, $F(e^{i\theta_0}) + F(e^{i\theta_0})^*$ is positive semidefinite for all $\theta_0 \in (0, \pi)$ that is not a pole of $G(z)$. Therefore, by continuity, we have $F(-1) + F(-1)^* \geq 0$.

(If). Let F be given by (11). Since $F(z)$ is symmetric, real, rational, proper, discrete-time positive real and $G(-1) = G(-1)^\top$, it is sufficient to show that

$$G_0(z) \stackrel{\text{def}}{=} \frac{z+1}{z-1} F(z)$$

is NI because $G_0(z)$ is NI if and only if $G(z) = G_0(z) + G(-1)$ is NI. We observe that $G_0(z)$ is proper, symmetric, real, rational, discrete-time and analytic in $|z| > 1$. Also, $F(z)$ and $G_0(z)$ have the same poles, with the possible exception of a pole at $z = 1$. Notice that $F(z)$ does not have a pole at $z = -1$ due to its construction in (11). Let $z_0 = e^{i\theta_0}$ with $\theta_0 \in (0, \pi)$. Assume z_0 is not a pole of $F(z)$. Then, it is not a pole of $G_0(z)$. We find

$$G_0(e^{i\theta_0}) = \frac{e^{i\theta_0} + 1}{e^{i\theta_0} - 1} F(e^{i\theta_0}) = -\frac{i \sin \theta_0}{1 - \cos \theta_0} F(e^{i\theta_0}),$$

so that

$$i[G_0(e^{i\theta_0}) - G_0(e^{i\theta_0})^*] = \frac{\sin \theta_0}{1 - \cos \theta_0} [F(e^{i\theta_0}) + F(e^{i\theta_0})^*] \geq 0,$$

because $F(e^{i\theta_0}) + F(e^{i\theta_0})^* \geq 0$. We now assume that $z_0 = e^{i\theta_0}$ with $\theta_0 \in (0, \pi)$ is a pole of $F(z)$. Then, it is also a pole of $G_0(z)$. Since $F(z)$ is PR, z_0 is a simple pole. Thus, z_0 is also a simple pole of $G_0(z)$. Moreover, the matrix $K_0 = e^{-i\theta_0} \lim_{z \rightarrow e^{i\theta_0}} (z - e^{i\theta_0}) F(z)$ is positive semidefinite, see Theorem 2.1. This then implies that

$$\begin{aligned} e^{-i\theta_0} \lim_{z \rightarrow z_0} (z - e^{i\theta_0}) i G_0(z) &= e^{-i\theta_0} \lim_{z \rightarrow z_0} i \frac{z+1}{z-1} (z - e^{i\theta_0}) F(z) \\ &= i \frac{e^{i\theta_0} + 1}{e^{i\theta_0} - 1} K_0 \\ &= \frac{\sin \theta_0}{1 - \cos \theta_0} K_0 \geq 0. \end{aligned}$$

When $z = 1$, $F(z)$ can either have no poles or a simple pole. Assume $z = 1$ is not a pole. Then, $G_0(z) = G_1(z) + \frac{2F(1)}{z-1}$ where $G_1(z)$ is analytic in a region near $z = 1$. Then, $K_0 = \lim_{z \rightarrow 1} (z-1) G_0(z) = 2F(1) = F(1) + F(1)^\top$ (due to $F(z)$ being symmetric), which is non-negative in view of Theorem 2.1.

Assume now that $z = 1$ is a simple pole of $F(z)$. We can write $F(z) = F_1(z) + \frac{A}{z-1}$, where $F_1(z)$ is analytic near $z = 1$ and $0 \leq A \leq 2F_1(1)$ (via Theorem 2.1, since $A \geq 0$ directly from the theorem statement and $0 \leq F(e^{i\theta}) + F(e^{i\theta})^* = F_1(e^{i\theta}) + F_1(e^{i\theta})^* - A$ implies $A \leq 2F_1(1)$ in the limit as $\theta \rightarrow 0$ due to continuity and $F_1(1)$ being symmetric).

Hence,

$$\begin{aligned} G_0(z) &= \frac{z+1}{z-1} F(z) = \frac{z+1}{z-1} F_1(z) + \frac{z+1}{(z-1)^2} A \\ &= G_2(z) + \frac{2F_1(1) + A}{z-1} + \frac{2A}{(z-1)^2}, \end{aligned}$$

where $G_2(z)$ is analytic in the neighbourhood of $z = 1$. Thus, the residue and the quadratic residue are $A_1 = A + 2F_1(1)$ and $A_2 = 2A$, and the condition that ensures that $F(z)$ is PR now guarantees that $A_2 \geq 0$ and $A_1 \geq A_2$, so that $G_0(z)$ is NI. \blacksquare

Lemma 3.8 Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper, NI transfer function with no poles at $z = -1$. Then

- $G(\infty) = G(\infty)^\top$;
- $G(-1)$ exists and $G(-1) = G(-1)^\top$.

Furthermore, let $\begin{bmatrix} A & B \\ C & D \end{bmatrix}$ be a minimal state-space realization of $G(z)$. Then,

- $C(I+A)^{-1}B = B^\top(I+A^\top)^{-1}C^\top$;
- $F(z) = \frac{z-1}{z+1} [G(z) - G(-1)]$ has a state-space realization

$$\left[\begin{array}{c|c} A & B \\ \hline C(A-I)(A+I)^{-1} & C(A+I)^{-1}B \end{array} \right]$$

which is minimal when A has no eigenvalues at 1.

Proof: Since $G(z)$ is symmetric, i.e. $G(z) = G(z)^\top$ for all $|z| > 1$, we obtain $G(\infty) = G(\infty)^\top$ via a limiting argument. Since $G(z)$ has no poles at $z = -1$, it follows that $G(-1)$ exists. Now, $G(z) = G(z)^\top$ for all $|z| > 1$ implies that $G(-1) = G(-1)^\top$ via continuity and a limiting argument.

From $G(-1) = G(-1)^\top$ and $D = D^\top$, it immediately follows that $C(I+A)^{-1}B = B^\top(I+A^\top)^{-1}C^\top$.

Let us now consider a state-space realization of $F(z)$. A realization of the transfer function matrix $\frac{z-1}{z+1}I$ is given by

$$\left[\begin{array}{c|c} -I & I \\ \hline -2I & I \end{array} \right], \text{ while a realization of the term } G(z) - G(-1) = C(zI - A)^{-1} + C(A+I)^{-1}B \text{ is given by } \left[\begin{array}{c|c} A & B \\ \hline C & C(A+I)^{-1}B \end{array} \right].$$

Thus, a realization for $F(z)$ is given by

$$\left[\begin{array}{cc|c} -I & C & C(A+I)^{-1}B \\ 0 & A & B \\ \hline -2I & C & C(A+I)^{-1}B \end{array} \right].$$

Changing state coordinates via

$$T = \begin{bmatrix} I & C(I+A)^{-1} \\ 0 & I \end{bmatrix}$$

yields

$$\begin{aligned} F(z) &= \left[\begin{array}{cc|c} -I & 0 & 0 \\ 0 & A & B \\ \hline -2I & C[I-2(I+A)^{-1}] & C(A+I)^{-1}B \end{array} \right] \\ &= \left[\begin{array}{cc|c} -I & 0 & 0 \\ 0 & A & B \\ \hline -2I & C(A-I)(I+A)^{-1} & C(A+I)^{-1}B \end{array} \right]. \end{aligned}$$

This realization is not minimal because it is easily seen that it is not completely reachable. Eliminating the non-reachable part one obtains

$$F(z) = \left[\begin{array}{c|c} A & B \\ \hline C(A-I)(A+I)^{-1} & C(A+I)^{-1}B \end{array} \right],$$

which is minimal if $\det(A-I) \neq 0$. \blacksquare

Remark 3.3 We have derived the condition $G(-1) = G(-1)^\top$ as a consequence of the symmetry of $G(z)$. However, if we consider, in the spirit of Remark 3.2, the possibly non-symmetric case, then condition $G(-1) = G(-1)^\top$ still holds. More precisely, assuming that rational NI systems are defined by conditions (i)-(v) of Lemma 3.2 (and that symmetry is not assumed), we have that if -1 is not a pole of $G(z)$ then $G(-1) = G(-1)^\top$. In fact, since by condition (ii) of Lemma 3.2, $i[G(e^{i\theta}) - G(e^{i\theta})^*] \geq 0$ for all $\theta \in (0, \pi)$ (except for the values of θ for which $z = e^{i\theta}$ is a pole of $G(z)$), we can use continuity and conclude that $i[G(-1) - G(-1)^*] \geq 0$, but $G(-1)$ is real so that $i[G(-1) - G(-1)^\top]$ is positive semi-definite. The diagonal entries of $i[G(-1) - G(-1)^\top]$ are zero so that we necessarily have $G(-1) - G(-1)^\top = 0$. Similarly, assuming that rational NI systems are defined by conditions (i)-(v) of Lemma 3.2 (and that symmetry is not assumed), we have that if 1 is not a pole of $G(z)$ then $G(1) = G(1)^\top$. In this non-symmetric setting, it is easy to check that the result analogue to Lemma 3.7 is that $G(z)$ without poles in -1 is NI if and only if $F(z)$ defined by (11) is PR and $G(-1) = G(-1)^\top$.

We are now in a position to give a discrete-time negative imaginary lemma that gives a complete state-space characterization of NI systems. Different grades of strength of continuous-time negative imaginary lemmas are given in [12,13,23].

Theorem 3.2 Let $\begin{bmatrix} A & B \\ C & D \end{bmatrix}$ be a minimal state-space realization of a discrete-time, symmetric, real, rational, proper transfer function $G(z)$. Suppose $\det(I+A) \neq 0$ and $\det(I-A) \neq 0$. Then, $G(z)$ is NI if and only if there exists a real matrix $X = X^\top > 0$ such that

$$X - A^\top X A \geq 0 \quad \text{and} \quad C = -B^\top (A^\top - I)^{-1} X (A+I). \quad (12)$$

Proof: First, note that

$$A(A-I)^{-1} = I + (A-I)^{-1}. \quad (13)$$

Now, in view of Lemma 3.7, $G(z)$ is NI if and only if $F(z) = \frac{z-1}{z+1} [G(z) - G(-1)]$ is PR. By Lemma 3.8, this is equivalent to

$$\left[\begin{array}{c|c} A & B \\ \hline C(A-I)(A+I)^{-1} & C(A+I)^{-1}B \end{array} \right]$$

being PR. Using Lemma 2.1, the latter conditions are equivalent to existence of $X > 0$ and L, W such that

$$X - A^\top X A = L^\top L, \quad (14)$$

$$(A^\top + I)^{-1} (A^\top - I) C^\top - A^\top X B = L^\top W, \quad (15)$$

$$C(A+I)^{-1} B + B^\top (A^\top + I)^{-1} C^\top - B^\top X B = W^\top W. \quad (16)$$

Eq. (15) can be written as

$$C = (W^\top L + B^\top X A) (A-I)^{-1} (A+I),$$

which can be substituted into (16) to give

$$\begin{aligned} B^\top X [I + (A-I)^{-1}] B + B^\top [I + (A^\top - I)^{-1}] X B - B^\top X B \\ = W^\top W - W^\top L (A-I)^{-1} B - B^\top (A^\top - I)^{-1} L^\top W \end{aligned}$$

in view of (13). This equation can also be written as

$$\begin{aligned} B^\top X (A-I)^{-1} B + B^\top (A^\top - I)^{-1} X B + B^\top X B \\ + B^\top (A^\top - I)^{-1} L^\top L (A-I)^{-1} B \\ = [W - L(A-I)^{-1} B]^\top [W - L(A-I)^{-1} B]. \end{aligned}$$

Substituting the term $L^\top L$ of (14) into the latter yields

$$\begin{aligned} B^\top X (A-I)^{-1} B + B^\top (A^\top - I)^{-1} X B + B^\top X B + \\ B^\top (A^\top - I)^{-1} X (A-I)^{-1} B - B^\top (A^\top - I)^{-1} A^\top X A (A-I)^{-1} B \\ = [W - L(A-I)^{-1} B]^\top [W - L(A-I)^{-1} B]. \end{aligned}$$

Using (13), it is easily seen that the left hand-side of this equation is equal to zero, so that $W = L(A-I)^{-1} B$. This means that $G(z)$ is NI if and only if there exists $X > 0$ such that $X - A^\top X A \geq 0$ and $C = [B^\top (A^\top - I)^{-1} (X - A^\top X A) + B^\top X A] (A-I)^{-1} (A+I)$.

Now, using (13), $G(z)$ is NI if and only if there exists $X > 0$ such that $X - A^\top X A \geq 0$ and $C(A+I)^{-1} = -B^\top(A^\top - I)^{-1}X$. ■

The conditions given in Theorem 3.2 and Corollary 3.1 are non-strict linear matrix inequalities in the Lyapunov variable X or Y which yield a set of convex conditions that can be solved via commercially available software. The conditions are necessary and sufficient. If one were to tighten the non-strict inequality to a strict inequality, a subset of negative imaginary systems would be obtained, see [13] for detailed discussions on this in the continuous-time setting.

Corollary 3.1 *Let the suppositions of Theorem 3.2 hold. Then $G(z)$ is NI if and only if there exists $Y = Y^\top > 0$ such that $Y - AY A^\top \geq 0$ and $B = -(A-I)Y(A^\top + I)^{-1}C^\top$.*

Proof: The result follows by letting $Y = X^{-1}$ and noting that $X - A^\top X A \geq 0$ is equivalent to $\begin{bmatrix} X & A^\top \\ A & X^{-1} \end{bmatrix} \geq 0$, which is in turn equivalent to $X^{-1} - AX^{-1}A^\top \geq 0$. ■

We next show that $G(1)$ and $G(-1)$ can be ordered for discrete-time negative imaginary systems.

Lemma 3.9 *Let $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper, NI (resp. WSNI) transfer function with no poles at $+1$ and -1 . Then*

$$G(1) - G(-1) \geq 0 \text{ (resp. } > 0 \text{)}.$$

Proof: Using Theorem 3.2 and a minimal realization for $G(z)$, we find

$$\begin{aligned} G(1) - G(-1) &= C(I-A)^{-1}B + D - C(-I-A)^{-1}B - D \\ &= C[(I-A)^{-1} + (I+A)^{-1}]B \\ &= 2C(I+A)^{-1}(I-A)^{-1}B \\ &= -2B^\top(A^\top - I)^{-1}X(I-A)^{-1}B \\ &= 2B^\top(I-A)^{-\top}X(I-A)^{-1}B \geq 0. \end{aligned}$$

This concludes the proof when G is NI.

Now, we focus on G being WSNI. The strict inequality result will be proven via a contra-positive argument. Suppose there exists an $x \in \mathbb{R}^m$ such that $[G(1) - G(-1)]x = 0$. Then $B^\top(I-A)^{-\top}X(I-A)^{-1}Bx = 0$ which implies that $Bx = 0$ as $X > 0$. This then implies that $G(e^{i\theta})x = Dx \ \forall \theta \in (0, \pi)$, i.e.,

$$(G(e^{i\theta}) - D)x = 0, \quad \forall \theta \in (0, \pi).$$

Since G is WSNI, $i[G(e^{i\theta}) - G(e^{i\theta})^*]$ is positive definite for all $\theta \in (0, \pi)$ so that if, for $\theta_0 \in (0, \pi)$, x is such

that $x^*[i(G(e^{i\theta_0}) - G(e^{i\theta_0})^*)]x = 0$, we can conclude that $x = 0$. Now recall that $D = D^\top$. Hence, $x^*[i(G(e^{i\theta_0}) - G(e^{i\theta_0})^*)]x = ix^*[(G(e^{i\theta_0}) - D) - (G(e^{i\theta_0}) - D)^*]x = 0$. Hence $x = 0$, so that $[G(1) - G(-1)]$ must be nonsingular. This completes the proof. ■

4 Feedback interconnections and internal stability

The following result shows under what circumstances are NI, WSNI and SSNI properties preserved when such systems are interconnected in feedback. Given complex matrices S_1, S_2 and complex vectors $y_1, y_2, u_1, u_2, \alpha, \beta$ of compatible dimension satisfying

$$\begin{bmatrix} y_1 \\ \alpha \end{bmatrix} = S_1 \begin{bmatrix} u_1 \\ \beta \end{bmatrix} \text{ and } \begin{bmatrix} \beta \\ y_2 \end{bmatrix} = S_2 \begin{bmatrix} \alpha \\ u_2 \end{bmatrix}, \text{ let}$$

$S_1 \star S_2$ denote the Redheffer star product which maps $\begin{bmatrix} u_1 \\ u_2 \end{bmatrix}$ to $\begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$. Furthermore, let $F_\ell(S_1, S_2^{(1,1)})$ (resp. $F_u(S_2, S_1^{(2,2)})$)

denote the lower (resp. upper) linear fractional transformation. Finally, let $[P, Q]$ denote the positive feedback interconnection between systems P and Q .

Lemma 4.1 *Let $S_1 : \mathbb{C} \rightarrow \mathbb{C}^{m_1 \times m_1}$ be NI (resp. WSNI or SSNI) and $S_2 : \mathbb{C} \rightarrow \mathbb{C}^{m_2 \times m_2}$ be NI (resp. WSNI or SSNI). Let $0 < a, b \leq \min\{m_1, m_2\}$ and suppose the feedback interconnection corresponding to the Redheffer Star product $S_1 \star S_2$ be internally stable.² Then $S_1 \star S_2$ is NI (resp. WSNI or SSNI).*

Furthermore, if

- $a = b = m_2 < m_1$, then $S_1 \star S_2 = F_\ell(S_1, S_2)$;
- $a = b = m_1 < m_2$, then $S_1 \star S_2 = F_u(S_2, S_1)$;
- $a = b = m_2 = m_1/2$, $S_1 = \begin{bmatrix} P & I_a \\ I_a & 0 \end{bmatrix}$ and $S_2 = Q$, then $S_1 \star S_2 = P + Q$;
- $a = b = m_2 = m_1/2$, $S_1 = \begin{bmatrix} 0 & I_a \\ I_a & P \end{bmatrix}$ and $S_2 = Q$, then $S_1 \star S_2 = Q(I_a - PQ)^{-1}$;
- $2a = 2b = m_1 = m_2$, $S_1 = \begin{bmatrix} 0 & I_a \\ I_a & P \end{bmatrix}$ and $S_2 = \begin{bmatrix} Q & I_a \\ I_a & 0 \end{bmatrix}$, then $S_1 \star S_2 = \begin{bmatrix} -P & I_a \\ I_a & -Q \end{bmatrix}^{-1} = \begin{bmatrix} Q(I_a - PQ)^{-1} & (I_a - QP)^{-1} \\ (I_a - PQ)^{-1} & P(I_a - QP)^{-1} \end{bmatrix}$ which corresponds to the positive feedback interconnection $[P, Q]$.

² This is the standard meaning of ‘‘internal stability’’, i.e. add two extra exogenous input signals to the internal signals and ensure that all output signals and all internal signals are energy-bounded for any energy-bounded exogenous input excitation.

Proof: Given $S_1(z), S_2(z)$ and complex vectors $y_1, y_2, u_1, u_2, \alpha, \beta$

of compatible dimension satisfying $\begin{bmatrix} y_1 \\ \alpha \end{bmatrix} = S_1(z) \begin{bmatrix} u_1 \\ \beta \end{bmatrix}$ and

$$\begin{bmatrix} \beta \\ y_2 \end{bmatrix} = S_2(z) \begin{bmatrix} \alpha \\ u_2 \end{bmatrix}, \text{ it follows that } \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = S_1(z) \star S_2(z) \begin{bmatrix} u_1 \\ u_2 \end{bmatrix}.$$

Then, for all $\begin{bmatrix} u_1 \\ \beta \end{bmatrix} \in \mathbb{C}^{m_1}, \begin{bmatrix} \alpha \\ u_2 \end{bmatrix} \in \mathbb{C}^{m_2}$:

$$\begin{aligned} & [u_1^* \ u_2^*] [i([S_1(z) \star S_2(z)] - [S_1(z) \star S_2(z)]^*)] \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} \\ &= i[u_1^* \ u_2^*] \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} - i[y_1^* \ y_2^*] \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} \\ &= i[u_1^* \ \beta^*] \begin{bmatrix} y_1 \\ \alpha \end{bmatrix} - i[y_1^* \ \alpha^*] \begin{bmatrix} u_1 \\ \beta \end{bmatrix} \\ &\quad + i[\alpha^* \ u_2^*] \begin{bmatrix} \beta \\ y_2 \end{bmatrix} - i[\beta^* \ y_2^*] \begin{bmatrix} \alpha \\ u_2 \end{bmatrix} \\ &= [u_1^* \ \beta^*] [i(S_1(z) - S_1(z)^*)] \begin{bmatrix} u_1 \\ \beta \end{bmatrix} \\ &\quad + [\alpha^* \ u_2^*] [i(S_2(z) - S_2(z)^*)] \begin{bmatrix} \alpha \\ u_2 \end{bmatrix}. \end{aligned}$$

Since the Redheffer star interconnection is internally stable, the three respective results (NI, WSNI, SSNI) then follow by applying Definition 3.1, Lemma 3.3 or Lemma 3.5 respectively on the corresponding domains of $z \in \mathbb{C}$ for $S_1(z)$ and $S_2(z)$.

The five cases where a, b, S_1 and S_2 are restricted are trivial consequences of a Redheffer calculation. ■

Example 4.1 This example shows that it is not possible to mix and match properties of S_1 and S_2 for the strict results in Lemma 4.1 to hold.

Let $S_1 = \begin{bmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} & \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ \begin{pmatrix} 1 & 0 \end{pmatrix} & 0 \end{bmatrix}$ which is clearly NI and let $S_2 = z^{-1}$ which is clearly SSNI (and hence also WSNI and hence also NI). Then $S_1 \star S_2 = \begin{bmatrix} 1+z^{-1} & 0 \\ 0 & 1 \end{bmatrix}$ which is only NI (and not WSNI nor SSNI).

The following stability theorem here applies to real, rational, proper systems but invokes only the interconnection of NI and WSNI systems. It is the discrete-time counterpart of Theorem 5 in [12].

Theorem 4.1 Let $P : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper, NI system with no poles at $+1$ and -1 , and

let $Q : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ be a discrete-time, real, rational, proper, WSNI system. Suppose $P(-1)Q(-1) = 0$ and $Q(-1) \geq 0$. Then

$$[P, Q] \text{ is internally stable} \Leftrightarrow \bar{\lambda}(P(1)Q(1)) < 1.$$

Proof: The proof trivially follows by applying [12, Theorem 5] or [29, Theorem 1] on the systems $M(s) = P(\frac{1+s}{1-s})$ and $N(s) = Q(\frac{1+s}{1-s})$ obtained through the bilinear transformation $z = \frac{1+s}{1-s}$. ■

Concluding remarks

In this paper we presented a definition of negative imaginary systems for discrete-time systems that hinges entirely on properties of the transfer function matrix and not on a real, rational, proper, finite-dimensional realization. We have drawn a full picture which illustrates the relationship that exists, in the discrete-time, between the notions of positive real and negative imaginary systems, as well as strictly positive real and strictly negative imaginary systems. Indeed, as it happens for the classical theory of positive real systems, even for negative imaginary systems our definitions can be viewed as a single definition referred to different analyticity domains. In fact, we can define a function $G : \mathbb{C} \rightarrow \mathbb{C}^{m \times m}$ analytic in the open set $\Omega = \{z \in \mathbb{C} : |z| > 1\}$, to be *skew-imaginary* if

- $i[G(z) - G(z)^*] \geq 0$ for all $z \in \Omega$ such that $\Im\{z\} > 0$;
- $i[G(z) - G(z)^*] = 0$ for all $z \in \Omega$ such that $\Im\{z\} = 0$;
- $i[G(z) - G(z)^*] \leq 0$ for all $z \in \Omega$ such that $\Im\{z\} < 0$.

Then, it is clear that a function is NI if it is analytic in Ω and skew-imaginary there.

Finally, we have derived stability analysis results for the interconnections of NI and WSNI systems.

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